8 Hensel's lemma la completions

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f Hensel's lemma

We will work with VK º7L.

We will figure out Hensel's lemma starting from a very basic question: given a valued field (K,N) and fEOVEX], how do we find a root of f in Or? A necessary condition is certainly that FEKV[X] has a root. Let ā EKV be such. This is of course not enough take revidues of the coefficients indeed, it may very well be that f(a) = 0. Yet, we may be optimistic and hope for bely such may  $\overline{b} = \overline{a}$  in f(b) = 0. How do find such a b? The neuran requirement that  $\overline{f}(\overline{a}) = \overline{O}$  is really the same as v(f(a)) 70, which insuitively speaking means that f(a) is very man - i.e., a is an approximation of a root of f. What we might be tempted to do is approach this like a numerical problem, so we might try to find a beller approximation of a roop than a, i.e.  $b \in O_V$  s, that v(f(b)) = v(f(a)). To do that, we will use Newton's method:  $b := a - \frac{f(a)}{f'(a)}$ F(X) Leb's ne whether b is really a better f(a) × approximation: f(b) $f(b) = f\left(a - \frac{f(a)}{f'(a)}\right) =$ 

$$Taylor expansion = f(a) + \left(-\frac{f(a)}{f'(a)}\right) f'(a) + \left(\frac{f(a)}{f'(a)}\right)^2 \cdot c$$
so,  

$$f(b) = f(a) - f(a) + \left(\frac{f(a)}{f'(a)}\right)^2 \cdot c = \left(\frac{f(a)}{f'(a)}\right)^1 \cdot c$$
hence  

$$v(f(b)) - v(f(a)) = 2v(f(a)) - 2v(f'(a)) + v(c) - v(f(a))$$

$$\Rightarrow v(f(a)) - 2v(f'(a))$$

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$$= f'(a) - f'$$

i) 
$$a_{n+1} = a_n - \frac{f(a_n)}{f'(a_n)}$$
  
ii)  $v(f(a_{n+1})) > v(f(a_n))$ .  
Indeed, we may see that  
 $v(a_2 \cdot a_n) = v\left(\frac{f(a_n)}{f'(a_n)}\right) = v\left(f'(a_n) \cdot \frac{f(a_n)}{f'(a_n)^2}\right)$   
 $= v\left(f'(a_n)\right) + v\left(f(a_n)\right) - 2v\left(f'(a_n)\right)$   
 $a_n = b = v(f'(b)) + v(f(b)) - 2v(f'(a))$   
 $v(f'(b)) = v(f'(a)) + v(f(b)) - 2v(f'(a))$   
 $= v(f'(a)) + v\left(\frac{f(a_n)}{f'(a_n)^2} \cdot c\right) - 2v(f'(a))$   
 $v(f'(a)) > 0 > 2v(f(a)) - 2v(f'(a)) + v(c) - 2v(f'(a))$   
 $v(c) > 0 > 2 \left[v(f(a)) - 2v(f'(a))\right]$   
and more generally,  $v(a_{n+n} - a_n) \ge 2^n \varepsilon$ , so  
 $f_{n \to \infty}$   
This should ving a bell:  $(a_{n+n} \cdot c_n) \le 2v(f'(a))$  is cauchy! Bub wall, we  
dom't have a metric, right?

& INTERNIDE

Given  $v: K^{\times} \rightarrow 7L$ , we may give K a metric by  $d_{v}(a,b) = \exp(-v(a-b)) \in \mathbb{R}^{7,0}$  is  $\exp(-\infty) = 0$ .

Ne can thus consider  $(K, d_v)$  as a method shall and Consider Cauchy sequences in it. To complete - prin not indended - our proof we will need  $(a_n)_{n\in\mathbb{N}}$  to converge, i.e.  $(K, d_v)$  to be complete as a metric prace.

→ if (X,d) is not complete, we can always embed (X,d) as the dense subspace of a complete metric space (X,d), which is unique up to inometry over X. It is tuilt this way: 1)  $\hat{X} = h(\operatorname{cancny sequences}^2/_{N}, \text{ where } (a_n)_n \sim (b_n)_n \Leftrightarrow \lim d(a_n, b_n) = 0,$ 2)  $\hat{d}([a_n], [b_n]) = \lim d(a_n, b_n).$ Think of  $\mathbb{Q}, \text{ with } \hat{d}(a_1, b) = 1 a - b1.$  Then  $\hat{\mathbb{Q}} \simeq \mathbb{R}^{?}$ If we take  $d_{Y}$  instead, i.e.  $d_{P}(a_1b) = p^{-V_{P}(a-b)}$ , then  $(\hat{\mathbb{Q}}, \hat{d}_{V_{P}}) = : (\mathbb{Q}_{P}, d_{P})^{?}$ is a friend we will meets often. We could write  $\mathbb{Q}_{P} = \{\sum_{n \geq N} a_{n}p^{n} \mid N \in TL, a_{n} \in h \circ_{1}1, \dots p - N\} \}$ and do computations try course.

Then, if we assume  $(K_{1}v)$  is complete (i.e.,  $(K_{1}dv)$  is complete as a metric space),  $(a_{1n})_{n\in\mathbb{N}}$  has a huit  $x\in\mathbb{Q}$  and we have

 $f(an) \rightarrow f(\alpha) \quad as f is continuous in the dy-topology$   $\Rightarrow \quad f(\alpha) = 0. \quad \text{Since} \quad v(f(an)) \rightarrow as \quad and \quad v(f(\alpha)) = hur v(f(au))$   $Moreover, \quad v(\alpha-a) \neq V(f'(a)) \neq 0 \quad and \quad so \quad \overline{\alpha} = \overline{\alpha}.$   $Summing \quad up,$ 

HENSEL'S LEMIMA. Suppose (K, v),  $vK \ge 7L$ , is a complete valued field. Then, for any  $f \in Q(X)$  and a O = V such that v(f(a)) > 2v(f'(a)), there is  $\forall \in Ov$ such that  $v(\alpha - a) > v(f'(a))$  is  $f(\alpha) = O$ .



ii) (KII is hemetian,

in.) every polynomial of the form

$$\frac{x^{n} + x^{n-1} + a_{n-1} x^{n-2} + \dots + a_{0}}{with a_{i} \in \mathbb{M}_{+}, 0 \leq i \leq n \cdot 2_{1}, was a 4200 in K.}$$

$$\frac{PROSE, we prove (i) = (i^{-}) \Rightarrow (i^{-}i^{-}) \Rightarrow (i^{-}i^{-}) = \overline{P(a)} + \overline{D}.$$

$$(i \Rightarrow in). Take (eq_{i}(x)_{+} a \in D_{v} with \overline{P(a)} - \overline{D}, \overline{P(a)} + \overline{D}.$$

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$$(i \Rightarrow in). Take point in the ave done; no around with the window of v to b. If n=4, one are done; no around the v in the window of the midple extension!) with  $\overline{a_{i}} = \overline{e(a_{i})} = \overline{E(a_{i})} = \overline{a} = \overline{a_{i}},$ 

$$(i = in). we have$$

$$\overline{f} = (x_{i} A) x^{n-4}$$

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that then if 
$$\xi \notin D_{i}$$
,  $O_{i}^{i} = O_{i} \wedge L = \delta_{i} (O_{i}) \wedge L \neq O_{i}^{1}$ . We can then  
apply Neak Approximation and find  $\beta \in R = O_{i}^{i} \dots \cap O_{n}^{i} \in L$   
such that  $\beta - A \in M_{n}^{i}$ ,  $\beta \in M_{n}^{i}$  for  $i \neq A$ . As  $O_{i}^{i} \neq O_{i}^{i}$ ,  $\beta \notin K$ : let  
 $f$  be its minimum only moutial over  $K_{i}$  gay  
 $f(X) = X^{k} + a_{k-n} X^{k-n} + \cdots + a_{n}$   
for none  $a_{i} \in K$   
Say  $p - \beta_{n}, \dots, \beta_{k}$  are the conjugates of  $\beta$  in  $N$ . For  $\delta \in G \cap D_{i}$ ,  
 $\beta \in \delta(m_{1})$  by definition, so  $\sigma(\beta) \in m_{1}$  for all  $\sigma \in C \cap D$ . As the  
 $\beta_{2}, \dots, \beta_{k}$  are exactive  $\delta(\beta_{n})$  for  $\delta \in G \cap D_{i}$ ,  $\beta_{i} \in M_{n}$ , for all  $i$ .  
And now  
 $f(X) = TT (X - \beta_{i})$   
satisfies  
 $a_{k-4} = -i\beta_{1} + \dots + \beta_{k} N \in (1 + m_{n}) \cap K = 1 + m_{n}$   
and  $a_{i} \in m_{1} \cap K = m_{n}$ .  
Thus  
 $g(X) = \frac{f(a_{k-n}X)}{(a_{k-1})^{k}} = \chi^{k} + \chi^{k-n} + a_{k-2}^{k-1} \chi^{k-1} + \cdots + a_{n}^{k} \in O_{i}[X]$   
(annot have  $a \neq m_{i}$ .